

§0. recall $P: X \rightarrow S$

inner fib \Leftarrow coCartesian fib \Leftarrow left fib \Leftarrow Kan fib

X_S is an ∞ -cat

$f: S \rightarrow T$
have $f! : X_S \rightarrow X_T$
up to canonical equi

X_S is a Kan complex

same over a point

slogan.

- $\text{coCart fib}/S \simeq \text{Fun}(S, \text{Cat}_{\infty})$
- $\text{Cart fib}/S \simeq \text{Fun}(S^{\text{op}}, \text{Cat}_{\infty})$
- $\text{left fib}/S \simeq \text{Fun}(S, \text{Kan})$
- $\text{right fib}/S \simeq \text{Fun}(S^{\text{op}}, \text{Kan})$

§1. symmetric monoidal category $(\mathcal{C}, \otimes, I)$

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Baez, some definitions everybody should know

<https://math.ucr.edu/home/baez/qg-fall2004/definitions.pdf>

Data: $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$

$I \in \text{ob}(\mathcal{C})$

natural isomorphisms $\left\{ \begin{array}{l} \alpha_X: I \otimes X \cong X \quad \text{(left) unit} \\ \beta_{X,Y}: X \otimes Y \cong Y \otimes X \quad \text{commutator} \\ \gamma_{X,Y,Z}: (X \otimes Y) \otimes Z \cong X \otimes (Y \otimes Z) \quad \text{associator} \end{array} \right.$

Requirement: coherent diagrams,

including pentagon and hexagon diagrams

Examples: $(\text{Vect}_k, \otimes, k)$

$(\text{Space}, \times, *)$

(Sp, \wedge, \S)

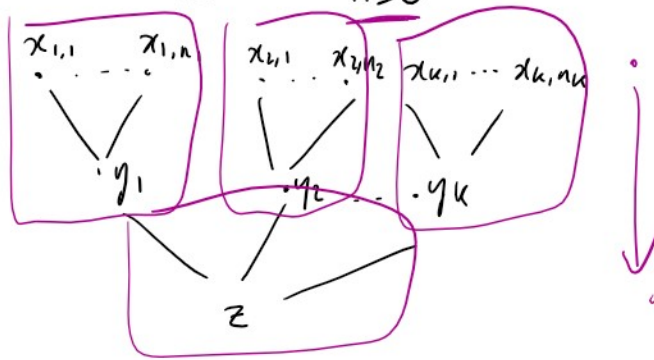
Multicategory \mathcal{M}

Objects. x, y

Multimorphisms. $\text{Mul}(x_1, \dots, x_n; y)$

$\sum_n \quad n \geq 0$

compositions



Example. $\text{Mul}(V_1, V_2; W) = \{ \text{bilinear maps } V_1 \times V_2 \rightarrow W \}$

example. $\text{Mult}(V_1, V_2; W) = \{ \text{bilinear maps } V_1 \times V_2 \rightarrow W \}$
 $= \{ \text{linear maps } V_1 \otimes V_2 \rightarrow W \}$

This is general:

s.m.c \longrightarrow multi-cat
 (e, \otimes, I) e^\otimes

What if we consider (x_1, \dots, x_n) as a single object?

multi-cat \longrightarrow categories / Fin_* .

$M \longmapsto P: M^\otimes \rightarrow \text{Fin}_*$
 M^\otimes obj (x_1, \dots, x_n) , $n \geq 0$, $x_i \in M$
 mor $(x_1, \dots, x_n) \rightarrow (y_1, \dots, y_m)$

consists of

$\left\{ \begin{array}{l} \alpha: \{0, 1, \dots, n\} \rightarrow \{0, 1, \dots, m\} \\ f_j \in \text{Mul}(X_{\alpha^{-1}(j)}; Y_j) \\ 1 \leq j \in m \end{array} \right.$

Fin_* : obj : $\langle n \rangle = \{0, 1, \dots, n\}$

mor : based maps

Why introduce the base point 0 ?

$\alpha: \langle n \rangle \rightarrow \langle m \rangle$
 can "forget" elements in $\langle n \rangle$.

Question:

• Important to forget points.

When and how to recover (e, \otimes, I) from $P: e^\otimes \rightarrow \text{Fin}_*$?

Answer:

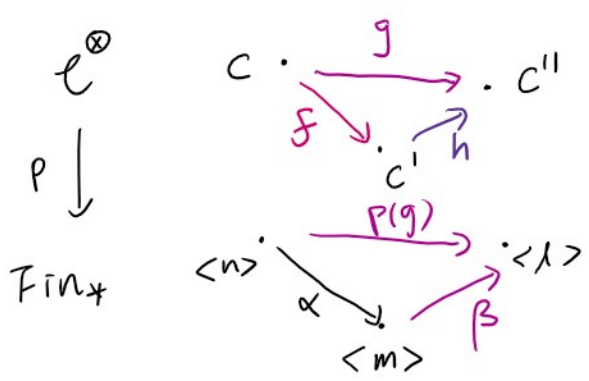
$\{ \text{s.m.c up to symmetric monoidal equivalence} \}$
 \Updownarrow
 $\{ p: \mathcal{D} \rightarrow \text{Fin}_* \text{ satisfying (M1) and (M2)} \}$

cocartesian condition
 segal condition

(M1) p is an op-fibration of categories

meaning that $\forall \alpha$, and c

$\exists f$, such that $\forall c''$, and
 g, β with $p(g) = \beta \circ \alpha$
 $\exists ! h$ with $p(h) = \beta$
 and $h \circ f = g$.



f is "p-cocartesian"

Content of (M1) for $p: \mathcal{D} \rightarrow \text{Fin}_*$

$\alpha: \langle n \rangle \rightarrow \langle m \rangle$ in Fin_*
filters over $\langle n \rangle$

$\Rightarrow \alpha_! : \mathcal{D}_{\langle n \rangle} \rightarrow \mathcal{D}_{\langle m \rangle}$ well defined up to
 Canonical isomorphism

Verification for $\mathcal{D} = e^\otimes$:

$e_{\langle n \rangle}^\otimes = e^n$ $\alpha: \langle n \rangle \rightarrow \langle m \rangle$
 $\alpha_!$ induced by $e^* \rightarrow e$
 $(x_1, \dots, x_k) \mapsto x_1 \otimes x_2 \otimes \dots \otimes x_k$

projections to 1
 injections

and $* \rightarrow \mathcal{C}$
 $* \mapsto I$

(M2) Segal condition

$$\begin{aligned} \rho^i : \langle n \rangle &\rightarrow \langle 1 \rangle \\ i &\mapsto 1 \\ \text{o.w. } &\mapsto 0 \end{aligned}$$

$$\mathcal{D}_{\langle n \rangle} \xrightarrow[\sim]{(P_1^1, \dots, P_1^n)} \mathcal{D}_{\langle 1 \rangle}^n$$

$$e^n = e^n$$

Verification for $\mathcal{D} = e^\otimes$:

$$\begin{aligned} \rho_1^i : e_{\langle n \rangle}^\otimes = e^n &\longrightarrow e_{\langle 1 \rangle}^\otimes = e \\ (x_1, \dots, x_n) &\mapsto x_i \end{aligned}$$

Content of (M2):

$$\mathcal{D}_{\langle n \rangle} \xrightarrow{\sim} \mathcal{D}_{\langle 1 \rangle}^n$$

not equal, but equivalent by built in data

To recover from $p: \mathcal{D} \rightarrow \text{Fin}_*$ the s.m.c. e

$$(M2) \Rightarrow \underline{e} = \underline{\mathcal{D}_{\langle 1 \rangle}}$$

$$\cdot \mathcal{D}_{\langle 0 \rangle} \simeq * \quad \underline{\langle 0 \rangle \rightarrow \langle 1 \rangle} \Rightarrow I \rightarrow e$$

$$\cdot \alpha : \langle 2 \rangle \rightarrow \langle 1 \rangle \quad 1, 2 \mapsto 1$$

$$\Rightarrow \underline{\otimes} : \mathcal{D}_{\langle 1 \rangle}^2 \xleftarrow[\text{(M2)}]{\sim} \mathcal{D}_{\langle 2 \rangle} \xrightarrow{\alpha!} \mathcal{D}_{\langle 1 \rangle}$$

$$\cdot \beta_{x,1} \text{ is given by } \beta : \langle 2 \rangle \rightarrow \langle 2 \rangle$$

.....

Summary

	classical	new
Data	<ul style="list-style-type: none"> \mathcal{C} \otimes, I α_x $\beta_{x,y}$ $\gamma_{x,y,z}$ 	$P: \mathcal{D} \rightarrow \text{Fin}_*$
Requirement	<ul style="list-style-type: none"> diagrams commute (e.g. MacLane pentagons) 	<ul style="list-style-type: none"> (M1) op-fibration \rightarrow <i>coCartesian fibration</i> (M2) segal

remark: active, inert

Advantage of the new perspective:

- ① simpler data
- ② easier to generalize to ∞ -categories

generalization of (M1)

Def. A symmetric monoidal ∞ -category is

a *coCartesian fibration* $P: \mathcal{C}^{\otimes} \rightarrow N(\text{Fin}_*)$

such that (*) P_i induced by P^i determines an

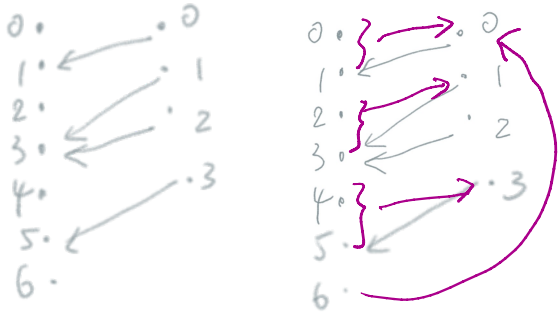
equivalence $\mathcal{C}_{\langle n \rangle}^{\otimes} \simeq (\mathcal{C}_{\langle 1 \rangle}^{\otimes})^n$ (M2)

Homework:

- ① Is $\text{id}: N(\text{Fin}_*) \rightarrow N(\text{Fin}_*)$ a symmetric monoidal ∞ -category?
- ② $\Delta^{\text{op}} \rightarrow \text{Fin}_*$ ~~Segal's~~ Γ -space *obsolete* in: $[m] \mapsto \langle m \rangle$ (category and cohomology theories)

$[m] \mapsto \langle m \rangle$ (categories and cohomology theories)

$([m] \leftarrow [n]) \mapsto (\langle m \rangle \rightarrow \langle n \rangle)$ (In fact, this is $S!$)



Is $N(S^!)$ a symmetric monoidal ∞ -category? An ∞ -operad? (No)

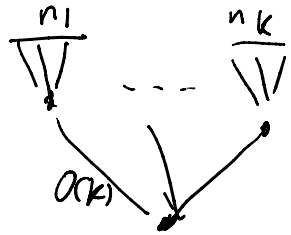
§2. Operads

data: sets $\mathcal{O}(n) \supset \Sigma_n$, $n \geq 0$

• $\eta: * \rightarrow \mathcal{O}(1)$ parameterizes $\text{id}: A \rightarrow A$

• $\gamma: \mathcal{O}(k) \times \mathcal{O}(n_1) \times \dots \times \mathcal{O}(n_k) \rightarrow \mathcal{O}(n_1 + \dots + n_k)$

requirement: unital, equivariance, associativity diagrams.



Quick observation:

• an operad = a multi-category with one object.

$\mathcal{O}(n) = \text{Mul}(\underbrace{*, \dots, *}_n; *)$ including $n=0$

\Rightarrow can form $p: \mathcal{O}^{\otimes} \rightarrow \text{Fin}_*$

$\left\{ \begin{array}{l} \text{obj} : \langle n \rangle = (*, \dots, *) \quad n \geq 0 \\ \text{mor} : \alpha : \langle m \rangle \rightarrow \langle n \rangle \end{array} \right.$

and $\left\{ x_j \in \mathcal{O}(n) \mid 1 \leq j \leq n \right\}$

Algebra

Examples:

Triv

$\text{Triv}(n) = \begin{cases} * & n=1 \\ \emptyset & n \neq 1 \end{cases}$

everything

Triv	$\{ \emptyset \} \quad n \neq 1$	everything
Triv [⊗]	obj: $\langle n \rangle$ mor: $\sigma \in \Sigma_n$	
E ₀	$E_0(n) = \begin{cases} * & n=0 \text{ or } 1 \\ \emptyset & \text{o.w.} \end{cases}$	based objects
E ₀ [⊗]	obj: $\langle n \rangle$ (based) mor: injections	
Comm	Comm(n) = *	comm monoids
Comm [⊗]	obj: $\langle n \rangle$ mor: based maps	Comm [⊗] Fin*
Ass	Ass(n) = Σ_n	(associative) monoids
Ass [⊗]	skipped	

Algebra over \mathcal{O} in $(\mathcal{C}, \otimes, I)$:

- $A \in \mathcal{C}$ with
- data: $\mathcal{O}(n) \times_{\Sigma_n} A^{\otimes n} \rightarrow A$
- requirement: unital and associativity diagrams

Content: $\mathcal{O}(n)$ parametrizes n -ary operations on A .
 If we allow $\mathcal{O}(n)$ to be spaces, we can parametrize
 "unique n -ary operations up to contractible choices".

Remark: We have seen the need for coherence data already
 in a symmetric monoidal category: there's
 $\int I: * \rightarrow \mathcal{C}$ look like $\begin{matrix} e^0 & \rightarrow & e^1 \\ \cdot 2 & & \cdot 1 \end{matrix}$

in a symmetric monoidal \mathcal{C}

$\left\{ \begin{array}{l} I: * \rightarrow \mathcal{C} \\ \otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C} \end{array} \right.$ look like $e^0 \rightarrow e^1$
 $e^2 \rightarrow e^1$

$\Rightarrow \mathcal{C}$ is some sort of algebra in $(\text{Cat}, \times, *)$

It's not a commutative algebra: we have coherence data $\alpha_x, \beta_{x,y}, \gamma_{x,y,z}$.

In a symmetric monoidal ∞ -category, even more coherence data is introduced.

strictness coherent data ↓ - +	cat	∞ -cat
	comm algebra	•
	symmetric monoidal category	•
	Σ_{∞} -algebra	symmetric monoidal ∞ -category comm algebra in s.m. ∞ -cat

Fact. $\text{CAlg}(\text{Cat}_{\infty}^{\times}) \simeq \text{Cat}_{\infty}^{\text{SMon}}$

$\text{Alg}(\text{Cat}_{\infty}^{\times}) \simeq \text{Cat}_{\infty}^{\text{Mon}}$

Homework.

③ \mathcal{C} : symmetric monoidal category. Find out the p -cocartesian edges in $p: \mathcal{C}^{\otimes} \rightarrow \text{Fin}_{\times}$

④ \mathcal{O} : operad. Do the same for $p: \mathcal{O}^{\otimes} \rightarrow \text{Fin}_{\times}$

§3. ∞ -operads.

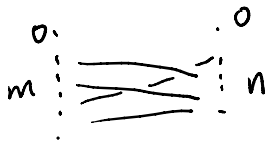
Recall

Def. A symmetric monoidal ∞ -category is a coCartesian fibration $p: \mathcal{C}^\otimes \rightarrow N(\text{Fin}_*)$

such that (*) p_i^i induced by p^i determines an equivalence $\mathcal{C}_{\langle n \rangle}^\otimes \simeq (\mathcal{C}_{\langle 1 \rangle}^\otimes)^n$.

Def. $\alpha: \langle m \rangle \rightarrow \langle n \rangle$ is inert if

$$|\alpha^{-1}\{i\}| = 1 \quad \underline{1 \leq i \leq n}$$



Example. $p^i: \langle n \rangle \rightarrow \langle 1 \rangle$ "coordinates"

Def. An ∞ -operad is $p: \mathcal{O}^\otimes \rightarrow N(\text{Fin}_*)$ between ∞ -categories such that

(1) \forall inert $\alpha: \langle m \rangle \rightarrow \langle n \rangle$, $c \in \mathcal{O}_{\langle m \rangle}^\otimes$, \exists p -coCartesian lift $\bar{\alpha}$.

(2) $\forall f: \langle m \rangle \rightarrow \langle n \rangle$, $c \in \mathcal{O}_{\langle m \rangle}^\otimes$, $c' \in \mathcal{O}_{\langle n \rangle}^\otimes$, choose p -coCartesian lifts of $p^i: \langle n \rangle \rightarrow \langle 1 \rangle$

$$\underline{p^i}: c' \rightarrow c'_i \in \mathcal{O}_{\langle 1 \rangle}^\otimes$$

Then

$\text{Map}_f \left(\dots, c'_i \right) \xrightarrow{(p^i \circ -)_{i=1}^n} \dots \xrightarrow{p^i \circ f} \dots$

component over f *component over $p^i \circ f$*

relax

$$\text{Map}_{\mathcal{O}^{\otimes}}^f(C, C') \xrightarrow{(\bar{p}^i \circ -)_{i=1}^n} \prod_{i=1}^n \text{Map}_{\mathcal{O}^{\otimes}}^{p^i \circ f}(C, C'_i)$$

is a homotopy equivalence.

(3) $\forall C_1, \dots, C_n \in \mathcal{O}_{\langle 1 \rangle}^{\otimes}$, $\exists C \in \mathcal{O}_{\langle n \rangle}^{\otimes}$ and p -cocartesian morphisms $\bar{p}^i: C \rightarrow C_i$ covering $p^i: \langle n \rangle \rightarrow \langle 1 \rangle$.

Rem (3') p^i induces $\phi: \mathcal{O}_{\langle n \rangle}^{\otimes} \xrightarrow{\sim} (\mathcal{O}_{\langle 1 \rangle}^{\otimes})^n \Rightarrow (3)$

(1) $\Rightarrow \phi$ well defined

segal condition

(2) $\Rightarrow \phi$ is fully faithful

(3) $\Rightarrow \phi$ is essentially surjective

What's (2)?

\mathcal{O} : (ordinary) operad before, $p: \mathcal{O}^{\otimes} \rightarrow \text{Fin}_*$

$\rightsquigarrow p: N(\mathcal{O}^{\otimes}) \rightarrow N(\text{Fin}_*)$ ∞ -operad

Verify (2): $f: \langle 4 \rangle \rightarrow \langle 2 \rangle \xrightarrow{p^i} \langle 1 \rangle$

$$\begin{array}{ccc} 0 \cdot & \longrightarrow & \cdot 0 \\ 1 \cdot & \longrightarrow & \cdot 1 \cdot 1 \\ 2 \cdot & \longrightarrow & \cdot 2 \cdot \\ 3 \cdot & \longrightarrow & \\ 4 \cdot & \longrightarrow & \end{array}$$

$$\text{Map}_{\mathcal{O}^{\otimes}}^f(4, 2) = \mathcal{O}(1) \times \mathcal{O}(2)$$

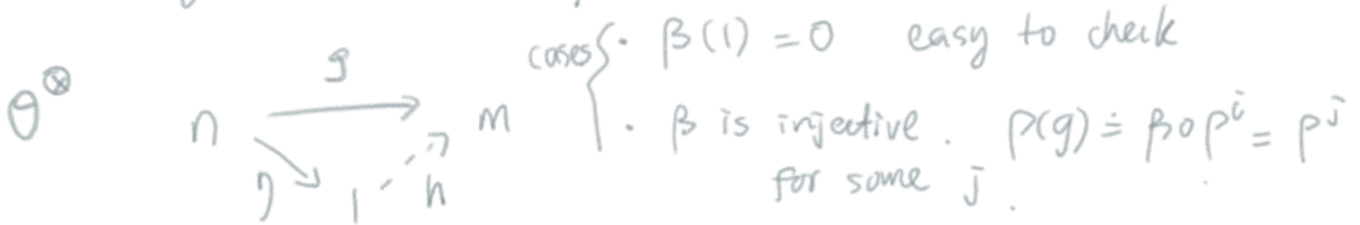
$$\downarrow \bar{p}^i \qquad \downarrow \text{projection}$$

$$\text{Map}_{\mathcal{O}^{\otimes}}^{p^i \circ f}(4, 1) = \begin{cases} \mathcal{O}(1) & i=1 \\ \mathcal{O}(2) & i=2 \end{cases}$$

Always pick $\bar{p}^i \in \mathcal{O}(1)$
 as $\eta \in \mathcal{O}(1)$.



* η is a comultiplication edge



$$g \leftrightarrow x \in \mathcal{O}(1)$$

$h \leftrightarrow$ element y in $\mathcal{O}(1)$
 want $y \circ \eta = x$
 identity in $\mathcal{O}(1)$
 monoid

Take $y = x$.

Then $\bar{p}^i \circ -$ is projection onto the component.

(2) is saying that morphisms in \mathcal{O}^\otimes all are made up by morphisms over $\langle n \rangle \rightarrow \langle 1 \rangle$.

Example. $\text{id}: N(\text{Fin}_*) \rightarrow N(\text{Fin}_*)$

\parallel
 Comm^\otimes

$\mathcal{O}(n)$

Example. $p: \mathcal{L}^\otimes \rightarrow N(\text{Fin}_*)$ a symmetric monoidal co-cat.

$\Rightarrow p$ is also an co-opelad.

— End —