


Intro to ω -category theory

Q: Why ω -category th'y?

A: Higher cat'l structures arise in algebra, geometry, math. physics, number th'y, etc.

Naive model: Take the (Quillen model) category of simplicially enriched categories, Cat_Δ .
(simplicial sets $\text{Set}_\Delta = \text{Fun}(\Delta^{op}, \text{set})$).

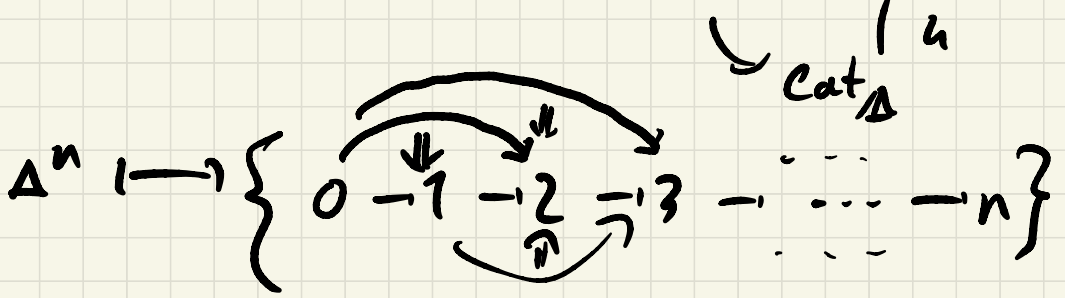
$$\mathcal{C} \in \text{Cat}_\Delta, \\ \mathcal{D} \vdash \text{Map}_{\mathcal{C}}(s, t) \xrightarrow{\sim} \text{Map}_{\mathcal{C}}(f(s), f(t)). \\ \forall s, t \in \mathcal{C}.$$

And $\forall t \in \mathcal{D}, \exists s \in \mathcal{C}$ and an equivalence $f(s) \xrightarrow{\sim} t$.

Not a good exponential in Cat_Δ .

Loosen this up. How?

$$\begin{array}{ccc} \text{Cat}^{\text{pt}} & \xrightarrow{\quad} & \text{Set}_\Delta, \\ \downarrow & \nearrow & (N\mathcal{C})_n = \text{Fun}(\begin{smallmatrix} [n] \\ \parallel \\ \{0 \rightarrow 1 \rightarrow \dots \rightarrow n\} \end{smallmatrix}, \mathcal{C}). \\ \text{Cat}_\Delta & \xrightarrow{N} & \mathcal{Q}[-] \triangleleft \text{Cat}_\Delta \end{array}$$



I.e. • $\text{Map}(0, 2) = \begin{pmatrix} 0 & 2 \\ \downarrow & \\ 0 & 12 \end{pmatrix} \cong \Delta^1$

• $\text{Map}(0, 3) = \begin{pmatrix} 0 & 3 & \rightarrow & 023 \\ \downarrow & & & \downarrow \\ 0 & 13 & \rightarrow & 0123 \end{pmatrix} \cong \Delta^1 \times \Delta^1$

etc.,

$$\text{Map}(i, j) \cong \begin{cases} \emptyset & j < i \\ \Delta^0 & j = i \\ (\Delta^1)^{j-i-1} & j > i \end{cases}$$

Thm: (Lurie) $N: \text{Cat}_\Delta \rightarrow \text{Set}_\Delta$

is right Quillen functor from the Bergner m.s. on Cat_Δ to Joyal m.s. on Set_Δ .

- fibrant objects are quasicategories $^{\text{Set}_\Delta}$.
- weak equivs are left and her. as before.

In Set_Δ , if \mathcal{C} and \mathcal{D} are ω -cats (Joyal fibrant) then $\text{Fun}(\mathcal{C}, \mathcal{D}) = \mathcal{D}^{\mathcal{C}}$ typically \rightarrow have

well-behaved.

(\star = join)

Tom: Define a monoidal product on Set_Δ
by $\Delta^m \star \Delta^n = \Delta^{m+1+n}$.

$$\{0 \rightarrow \dots \rightarrow m\} \star \{0 \rightarrow \dots \rightarrow n\} = \{0 \rightarrow 1 \rightarrow \dots \rightarrow m \rightarrow m+1 \rightarrow \dots \rightarrow m+n\}$$

Extend to Set_Δ by colims:

$$X \star Y = \text{colim}_{\substack{\Delta^m \rightarrow X, \\ \Delta^n \rightarrow Y}} \Delta^m \star \Delta^n.$$

If $f: \mathcal{D} \rightarrow \mathcal{C}$ is a functor, then

\mathcal{C}/f (the slice of \mathcal{C} over $f: \mathcal{D} \rightarrow \mathcal{C}$)

has n -simplices $(\mathcal{C}/f)_n = \{ \Delta^n \star \mathcal{D} \rightarrow \mathcal{C} \}$

$$\begin{array}{ccc} & \uparrow & \nearrow \\ & \phi \star \mathcal{D} \cong \mathcal{D}^{\uparrow} & f \end{array}$$

In particular, if $\mathcal{D} = *$,

$f \in \mathcal{C}$.

\mathcal{C}/f is just the usual "slice of \mathcal{C} over the object f ".

$$\{ \Delta^n \rightarrow \mathcal{C}/f \} = \{ \Delta^{n+1} \rightarrow \mathcal{C} \mid \text{final vertex is } f \}$$

Defn: If \mathcal{C} is an ω -cat, and $s, t \in \mathcal{C}$,
then $\text{Map}_{\mathcal{C}}(s, t) \xrightarrow{\text{fib}} \text{Fun}(\Delta^1, \mathcal{C}) \rightarrow \text{Fun}(\partial\Delta^1, \mathcal{C})$
is fiber over $(s, t) \in \mathcal{C} \times \mathcal{C}$.

Def'n: An object $t \in \mathcal{C}$ is terminal if $\forall s \in \mathcal{C}, \text{Map}_{\mathcal{C}}(s, t) \xrightarrow{\sim} *$.

Def'n: Let $f: \mathcal{D} \rightarrow \mathcal{C}$ be a functor of ω -cats. Then a limit of f is a terminal object of \mathcal{C}/f .

Eg: If \mathcal{C} is a fibrant simplicial category, then homotopy (co) limits in \mathcal{C} compute (co) limits in $N\mathcal{C}$.

Fibrations

If $S \in \text{Set}, \text{Set}/S \xleftarrow{\sim} \text{Fun}(S^{op}, \text{Set})$

Grothendieck construction $X = \int f \leftarrow f \rightarrow *$
 $\{ (s \in S, x \in f(s)) \}$

$X \xrightarrow{\Gamma} S$
 $\int f \xrightarrow{\Gamma} \int *$
 $\text{Cat}_{\infty} := N(\text{Cat})$
 $\text{Gpd}_{\infty} := N(\text{Kan})$

the simp. set of maps \mathcal{C} to \mathcal{D} is max'l Kan complex inside $\text{Cat} \subseteq \text{Set}_{\Delta}$
 $\text{Fun}(\mathcal{C}, \mathcal{D})$
 $\text{Kan} \subseteq \text{Set}_{\Delta}$

Thm: (Classifying spaces)

$$\text{Cat}_{\omega/S} \supset (\text{Cat}_{\omega/S})^{\text{right}} \xleftarrow{\sim} \text{Fun}(S^{\text{op}}, \text{Gpd}_{\omega})$$

$$\downarrow \int$$

$$\text{Gpd}_{\omega/S} \xleftarrow{\sim} \int f$$

$$\{(s \in S, x \in f(s))\}$$

"right fibrations"
"Cartesian fibrations".

This fails if S is not an ∞ -gpd:

(Lurie)

$$\text{Thm: } (\text{Cat}_{\omega/S})^{\text{cart}} \xleftarrow{\sim} \text{Fun}(S^{\text{op}}, \text{Cat}_{\omega})$$

$$\downarrow \int$$

$$\int f \xleftarrow{\sim} \int f$$

Moral: $p: X \rightarrow S$ is a cartesian fibration (contravariantly)

$\Leftrightarrow s \mapsto X_s$ is functorial in S .

Def'n: A map $p: X \rightarrow S$ in Cat_{ω} is a cartesian fibration if, for every map

$f: s \rightarrow t \in S$ and every object $y \in X$,

there exists a cartesian lift of f to y :

i.e. a map $g: x \rightarrow y \in X$ st $p(g) = f$

and $X_{t/g} \xrightarrow{\sim} X_{t/y} \times_{S_{t/e}} S_{t/f}$.

Ex. Vect \rightarrow Man, a cartesian fibration because vector bundles pull back.

There's a universal cocartesian fibration

$$\text{Cat}_{\infty}^{*//} \simeq \int \text{id}_{\text{Cat}_{\infty}} \longrightarrow \int^* \text{Cat}_{\infty} \simeq \text{Cat}_{\infty}$$

Because of this and $\text{Cat}_{\infty/S}^{\text{cocart}} \simeq \text{Fun}(S, \text{Cat}_{\infty})$ any cocartesian fibration pulls back:

$$\begin{array}{ccc} X & \longrightarrow & \text{Cat}_{\infty}^{*//} \\ \downarrow P & & \downarrow \\ S & \xrightarrow{f} & \text{Cat}_{\infty} \end{array}$$

Def'n: An adjunction in Cat_{∞} is a map $\mathcal{C} \rightarrow \mathcal{D}^1$ which is cartesian and cocartesian.

$$\begin{array}{ccc} \mathcal{C} = \mathcal{C}_0 & \begin{array}{c} \xrightarrow{f} \\ \xleftarrow{g} \end{array} & \mathcal{D} = \mathcal{D}_1 \end{array}$$

Want to define a subcategory of CAT_{∞} the presentable ω -cats.

Thm: (Lurie) If S is an ω -cat, then $\mathcal{D}(S) = \text{Fun}(S^{\text{op}}, \mathcal{G}pd_{\omega})$ is the universal cocompletion of S .

I.e. if T cocomplete, then

$$\text{Fun}^{\text{colim}}(\mathcal{D}(S), T) \xrightarrow{\sim} \text{Fun}(S; T).$$

restriction along Yoneda.

Def'n: An ω -cat \mathcal{C} is presentable if there is a small ω -cat S and a left adjoint $\mathcal{D}(S) \rightarrow \mathcal{C}$ ("small gens") st

- the right adjoint $\mathcal{C} \xrightarrow{R} \mathcal{D}(S)$ is fully faithful and R preserves k -filtered colims for some $k \gg 0$.

("small relative")

Thm: (Adjoint functor thm)

A functor of pres. ω -cats $f: \mathcal{C} \rightarrow \mathcal{D}$ is a left adjoint iff f preserves colims.